

Translate each to English, where the domain consists of all real numbers.

1. $\forall x \exists y (x < y)$

2. $\forall x \forall y ((x \geq 0 \wedge y \geq 0) \rightarrow xy \geq 0)$

Express so that negation immediately precedes predicate.

3. $\neg \forall x \exists y P(x, y)$

4. $\neg \forall x \forall y (P(x) \wedge F(y))$

What rule of inference is used in the following?

4. Alice is a math major. Therefore, Alice is a math or CS major.

5. If it snows today then school will close. School is not closed. Therefore it did not snow today.

TABLE 1 Rules of Inference.

<i>Rule of Inference</i>	<i>Tautology</i>	<i>Name</i>
$\begin{array}{l} p \\ p \rightarrow q \\ \hline \therefore q \end{array}$	$[p \wedge (p \rightarrow q)] \rightarrow q$	Modus ponens
$\begin{array}{l} \neg q \\ p \rightarrow q \\ \hline \therefore \neg p \end{array}$	$[\neg q \wedge (p \rightarrow q)] \rightarrow \neg p$	Modus tollens
$\begin{array}{l} p \rightarrow q \\ q \rightarrow r \\ \hline \therefore p \rightarrow r \end{array}$	$[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r)$	Hypothetical syllogism
$\begin{array}{l} p \vee q \\ \neg p \\ \hline \therefore q \end{array}$	$[(p \vee q) \wedge \neg p] \rightarrow q$	Disjunctive syllogism
$\begin{array}{l} p \\ \hline \therefore p \vee q \end{array}$	$p \rightarrow (p \vee q)$	Addition
$\begin{array}{l} p \wedge q \\ \hline \therefore p \end{array}$	$(p \wedge q) \rightarrow p$	Simplification
$\begin{array}{l} p \\ q \\ \hline \therefore p \wedge q \end{array}$	$[(p) \wedge (q)] \rightarrow (p \wedge q)$	Conjunction
$\begin{array}{l} p \vee q \\ \neg p \vee r \\ \hline \therefore q \vee r \end{array}$	$[(p \vee q) \wedge (\neg p \vee r)] \rightarrow (q \vee r)$	Resolution

TABLE 2 Rules of Inference for Quantified Statements.

<i>Rule of Inference</i>	<i>Name</i>
$\frac{\forall x P(x)}{\therefore P(c)}$	Universal instantiation
$\frac{P(c) \text{ for an arbitrary } c}{\therefore \forall x P(x)}$	Universal generalization
$\frac{\exists x P(x)}{\therefore P(c) \text{ for some element } c}$	Existential instantiation
$\frac{P(c) \text{ for some element } c}{\therefore \exists x P(x)}$	Existential generalization

For $x = -1, 0, 1$.

6. Given $\forall x P(x)$
What is $P(-1)$?

7. Given $P(c)$ for any c
What is $\forall x P(x)$?

Definition of even integer: $n = 2k$

Definition of odd integer: $n = 2k+1$

The sum of two even integers is even.

8. Restate as $p \rightarrow q$.

9. What is the hypothesis, p ?

10. What is the conclusion, q ?

11. Give a direct proof.

Definition of even integer: $n = 2k$

Definition of odd integer: $n = 2k+1$

If n is an even number then n^2 is an even number.

12. What is the hypothesis, p ?

13. What is the conclusion, q ?

14. What is the contrapositive?

15. Give a direct proof.

If n^2 is an odd integer then n is an odd integer.

16. Prove by contradiction.

Translate each to English, where the domain consists of all real numbers.

1. $\forall x \exists y (x < y)$ **For all x there exists a y such that $x < y$**
2. $\forall x \forall y ((x \geq 0 \wedge y \geq 0) \rightarrow xy \geq 0)$ **For all x and for all y, if $x \geq 0 \wedge y \geq 0$ then $xy \geq 0$**

Express so that negation immediately precedes predicate.

3. $\neg \forall x \exists y P(x, y) = \exists x \neg \exists y P(x, y) = \exists x \forall y \neg P(x, y)$
4. $\neg \forall x \forall y (P(x) \wedge F(y)) = \exists x \neg \forall y (P(x) \wedge F(y)) = \exists x \exists y \neg (P(x) \wedge F(y)) = \exists x \exists y (\neg P(x) \vee \neg F(y))$

What rule of inference is used in the following?

4. Alice is a math major. Therefore, Alice is a math or CS major.

Addition

5. If it snows today then school will close. School is not closed. Therefore it did not snow today.

Modus tollens

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For $x = -1, 0, 1$.

6. Given $\forall x P(x)$
What is $P(-1)$

True

7. Given $P(c)$ for
any c

What is $\forall x P(x)$

True

The sum of two even integers is even.

8. Restate as $p \rightarrow q$.

If x and y are even integers, then $x+y$ is even.

9. What is the hypothesis, p ?

10. What is the conclusion, q ?

11. Give a direct proof.

Prove by direct proof

"If x and y are even integers, then $x+y$ is even"

$\forall n (P(n) \rightarrow Q(n))$

Hypothesis

" x and y are even integers"

Show (conclusion) is true

" $x+y$ is even."

$P(n)$ Assume hypothesis, x and y even:
 $x = 2i, y=2j$

$x+y = 2i+2j$ Definition of even x and y

$= 2(i+j)$ Definition of even

\therefore even $x+y$

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Definition of odd integer: $n = 2k+1$

If n is an even number then n^2 is an even number.

12. What is the hypothesis, p ?

13. What is the conclusion, q ?

14. What is the contrapositive?

If n^2 is an odd number then
 n is an odd number .

15. Give a direct proof.

Prove by direct proof

"If n is an even integer, then n^2 is even"

$$\forall n (P(n) \rightarrow Q(n))$$

Hypothesis

" n is an even integer"

Show (conclusion) is true

"Given any even integer n , n^2 is even."

$P(n)$ Assume hypothesis, n is even: $n = 2k$

$n = 2k$ Definition of even n

$n^2 = (2k)^2$ Square both sides

$(2k)^2 = 4k^2$ Multiply

$4k^2 = 2(2k^2)$ Factor

$n^2 = 2(2k^2)$ Definition of even integer, $2k$.

$\therefore n^2$ is even

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Definition of even integer: $n = 2k$
Definition of odd integer: $n = 2k+1$

If n^2 is an odd integer then n is an odd integer.

16. Prove by contradiction.

Prove by contradiction
"If n is an odd integer, then n^2 is odd"
 $\forall n (P(n) \rightarrow Q(n))$
Hypothesis, p : " n^2 is an odd integer"
Conclusion, q : " n is odd integer."
 $\neg q$: n is even integer.

$n = 2k$	Assume $\neg q$, definition of even n
$n^2 = (2k)^2$	Square both sides
$(2k)^2 = 4k^2$	Multiply
$4k^2 = 2(2k^2)$	Factor
$n^2 = 2(2k^2)$	Definition of even integer, $2k$, contradicts hypothesis, n^2 is odd.
$\therefore n$ is odd	